KCL: A Declarative Language for Large-scale Configuration and Policy Management

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Background

01

Background



Configurations and Policies are important and crucial



Time	Event
2021.07	The Bilibili website in China went down because SLB Lua configuration code fell into an infinite loop with calculation errors
2021.10	KT Company in South Korea suffers major network interruption nationwide due to wrong routing configuration

Why Design KCL



Growing importance of modeling, constraint and scalability

Dynamic argument input

Helm

Pro	os.	Pro	OS.
•	Easy to write and read	•	Simple config logic su
•	Rich multi-language API	•	Dynamic argument in
•	Various Path Tools		
Co	ns.	Co	ns.
•	Redundant information	•	Increase of argument
•	Insufficient functionality e.g	g. it d	ifficult to maintain
abs	straction, constraint,	•	Insufficient functional
Tec	ch.	abs	straction, constraint,
•	JSON	Teo	ch.
•	YAML	•	Velocity
Product		•	Go Template
•	Kustomize	Pro	oduct
			** 1

Pros. Required programming features • Simple config logic support Code modularity Templates & Data abstraction Cons. Insufficient type constraints Increase of argument makes Insufficient restraint ability Insufficient functionality e.g. • Runtime error Tech. GCL **HCL JSONNET Product** Terraform

Pros. Rich config constraint syntax Unified type & value constraint • Configuration conflict checking writing with rich merge strategies Cons. Difficult to configuration override for multi-environment scenarios Runtime checks and limited performance Tech. CUE **Product Product** KusionStack KubeVela • ...

Pros. Model-centric & constraint-centric Scalability on separated block Static type system & analysis High Performance Cons. Expansion of different models requires investment in R&D Tech.

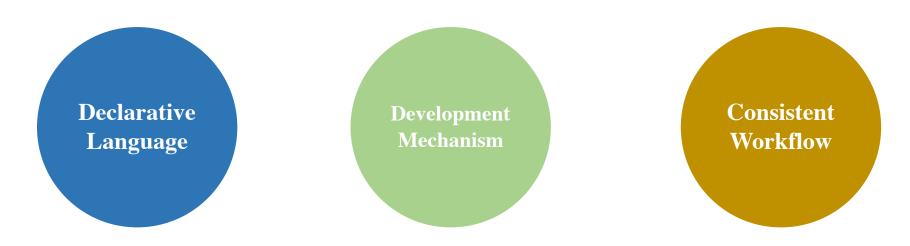
KCL

Modeled KV

Contribution



- Proposing the KCL declarative language, development mechanism, and consistent workflow to improve the large-scale efficiency and liberate multi-team collaborative productivity of operational development and operation systematically while ensuring stability for large-scale configuration and policy management.
- To date, the KCL has been used in more than 800 projects, and the average configuration writing and distributing time is shortened from more than 25 days to 2 days.





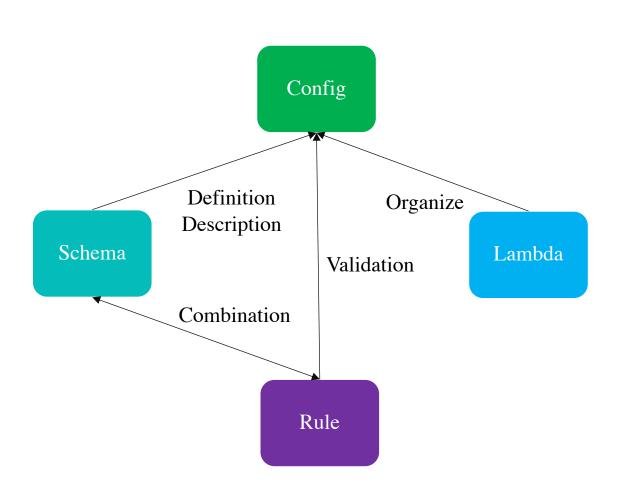
Design

02

Overview



KCL is an open source constraint-based record & functional language mainly used in configuration and policy scenarios



```
import units
     typa UnitTypa - unita NumberMultiplier
            import kubernetes.core.v1
 5
     SC
            deployment = v1.Deployment 
 6
                metadata.name = "nginx-deployment"
                metadata.labels.app = "nginx"
                spec = {
 9
                    replicas = 3
10
                    selector.matchLabels.app = "nginx"
        9
                   mixin OverQuotaMixin:
12
       10
                       overQuotaToleration = Toleration {
13
       11
                          key = "sigma.ali/is-over-quota"
     -12
                          operator = "Equal"
       13
                          value = "true"
       14
                          effect = "NoSchedule"
       15
                       } if data.overQuota else Undefined
       16
                       overQuotaMatchExpression = NodeSelectorRequirement {
       17
                          key = "sigma.ali/is-over-quota"
       18
               11
                          operator = "In"
               12
                          values = ["true"]
               13
                       } if data.overQuota else Undefined
```

Grammar



```
KCL := stmt^*
            stmt ::= schema \mid assert \mid import \mid assign \mid \dots
         schema := [\mathbf{schema}] id[`('parent\_id')'] [`['arguments']']` : 'schema\_body
   schema\_body ::= [mixin\_stmt] [schema\_context]
                       [check\_block]
    mixin\_stmt := mixin '['id^*']'
schema\_context ::= [attr|if \mid assign \mid assert \mid expr]^*
             attr := id[??] \cdot : Type [' = expr]
     check\_block := \mathbf{check} \cdot : \cdot check\_stmt^*
     check\_stmt ::= expr[`, 'msq]
          assign := id ' = ' expr
          assert := expr[`, `msg]
```

Fig. 1: (Subset of) KCL grammar

Core Pattern k op(T) v,

Type

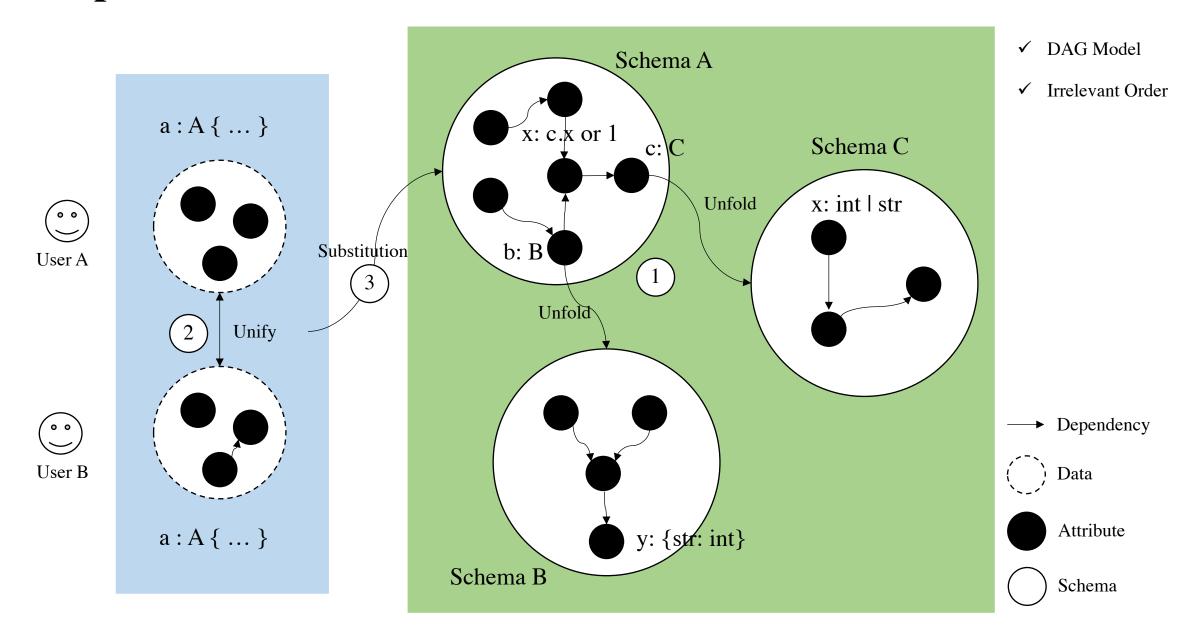


Table 1: KCL type notions

Type notion	Notation							
Boolean	$\frac{\Gamma \vdash \Diamond}{\Gamma \vdash boolean}$							
Integer	$\frac{\Gamma \vdash \Diamond}{\Gamma \vdash integer}$							
Float	$\frac{\Gamma \vdash \Diamond}{\Gamma \vdash float}$							
String	$\frac{\Gamma \vdash \Diamond}{\Gamma \vdash string}$							
Literal	$\frac{c \in \{boolean, integer, float, string\}}{\Gamma \vdash literalof(c)}$							
List	$\frac{\Gamma \vdash T \ T \neq Void}{\Gamma \vdash listof(T)}$							
Dict	$\Gamma \vdash T_1 \Gamma \vdash T_2 T_1 \neq \mathring{V} \circ id T_2 \neq V \circ id$							
Structure	$\frac{\Gamma \vdash dictof(T_k = T_1, T_v = T_2)}{\Gamma \vdash T_1 \dots \Gamma \vdash T_n T_i \neq V oid K_i \neq K_j, \text{when } i \neq j}$ $\Gamma \vdash structof(K_1 : T_1, \dots, K_n : T_n)$							
Union	$\frac{\Gamma \vdash T_1 \dots \Gamma \vdash T_n T_i \neq Void}{\Gamma \vdash unionof(T_1, \dots, T_n)}$							
Void	$\Gamma \vdash \Diamond$							
Any	$\frac{\Gamma \vdash Void}{\Gamma \vdash \diamondsuit}$							
Nothing	$rac{\Gamma dash Any}{\Gamma dash \lozenge} \over \Gamma dash Nothing}$							

Well-formed type $\Gamma \vdash \Diamond$.

Compilation



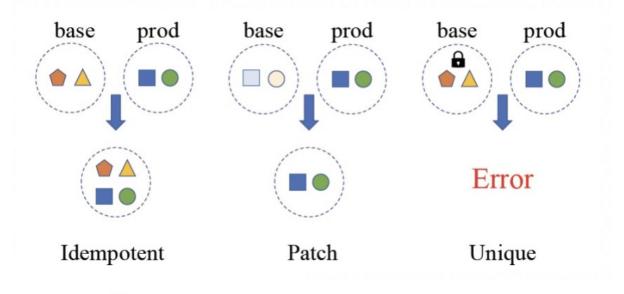
Configuration Substitution

```
Algorithm 1: The configuration substitution process.
   Input: the input configuration C_i
   Output: the merged configuration C_o
 1 C_o \leftarrow \{\};
 2 cache \leftarrow \{\};
 3 foreach k, e in C_i do
       if k in cache then
           C_o[k] = cache[k];
       else
           v = lookup(k, e);
           cache[k] = v;
       end
10 end
11
12 function lookup(k, e)
13 v \leftarrow \varnothing;
14 foreach k_d in dependence (e) do
       if k_d in cache then
15
           v = cache[k_d];
16
       else
17
           v = \text{lookup}(k_d, C_i[k_d]);
18
19
       end
20 end
21 return v;
22 end
```

Main Steps

- 1. Traverse all the key k and expression e of the input configuration C_i , and use the lookup function to substitute its dependent value recursively.
- 2. Store the calculated value in the output configuration and cache, which is used to avoid multiple calculations.
- 3. When all the key calculations are completed, we get the substitution completed configuration C_o .

Configuration Merge



Pattern
$$\sigma = \sum_{i=1}^{N} \{k_i, v_i, o_i\},$$

Merge
$$\sigma_u \cup \sigma_v = \sum_{j=1}^{N_u + N_v} \{k_j, v_j, o_j\}.$$

Result
$$\sum_{k=1}^{M} \{k_s, v_k, o_k\} = \{k_s, \{v_1, o_1\} \oplus \{v_2, o_2\} \oplus ... \oplus \{v_M, o_M\}\},$$

Configuration Properties

$$\sum_{k=1}^{M} \{k_s, v_k, o_k\} = \{k_s, \{v_1, o_1\} \oplus \{v_2, o_2\} \oplus \dots \oplus \{v_M, o_M\}\},\$$

where \oplus denotes the entry value merge operator.

Properties

- When the operation o_1 is the unique configuration operation, the calculation is invalid.
- When the operation o_1 is not the unique configuration operation, the configuration entry on the right side of the \oplus operator has a higher priority, so we can get $\{v_1, o_1\} \oplus \{v_2, o_2\} = \{v_1 \oplus v_2, o_2\}$
- When o_2 is an idempotent merge operation, there is a commutative law $\{v_2 \oplus v_1, o_2\} = \{v_1 \oplus v_2, o_2\}$ [37], and when the recursive partial order relationship $v_2 \subseteq v_1$ is satisfied, the calculation is valid. Besides, the partial

Configuration Properties

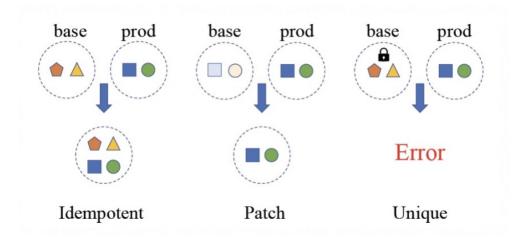
order calculation of two values can be defined as

$$v_2 \subseteq v_1 \Leftrightarrow \begin{cases} v_2 == v_1 \ if \ v_1 \in \{Int, Float, Boolean, String\} \\ \operatorname{len}(v_2) == \operatorname{len}(v_1) and \ \forall \ i, \ \operatorname{item}(v_2, i) \subseteq \operatorname{item}(v_1, i) \\ if \ v_1 \in List \ and \ if \ v_2 \in List \\ \forall \ k, \operatorname{item}(v_2, k) \subseteq \operatorname{item}(v_1, k) \\ if \ v_1 \in \{Dict, Structure\} \ and \\ if \ v_2 \in \{Dict, Structure\} \end{cases}$$

where i denotes the index of the list element, len denotes the length of the list element, k denotes the key of the dict and structure element, and item denotes the item of the list, dict and structure element.

- When the operation o_2 is an overwrite operation, $\{v_1 \oplus v_2, o_2\} = \{v_2, o_2\}$.
- When the operation o_2 is an append operation, it will try to add v_2 to the list of v_1 , and stop the calculation when v_1 is not a list type.

Example



• base.k

• prod.k

```
appConfiguration: frontend.Server {
    schedulingStrategy.resource = res.Resource {
        cpu = 100m
        memory = 100Mi
        disk = 1Gi
    }
}
```

Equivalent code



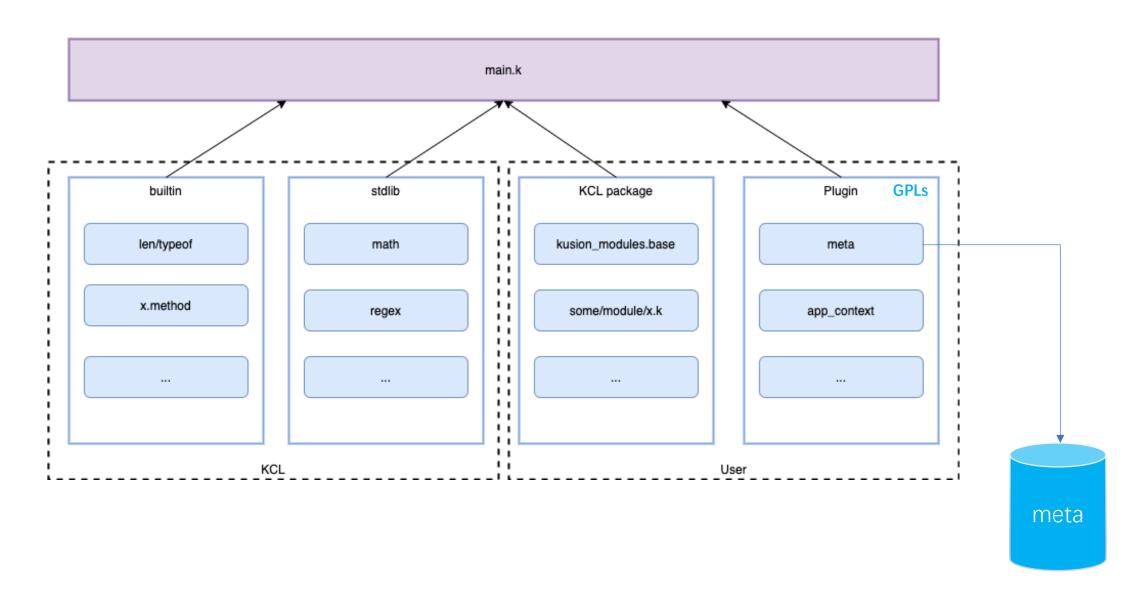
Automation

kcl -O appConfiguration.image="nginx:1.7.9"

```
import base.pkg.kusion_models.kube.frontend
    import base.pkg.kusion_models.kube.frontend.service
    import base.pkg.kusion_models.kube.frontend.container
    import base.pkg.kusion models.kube.templates.resource as res tpl
    # Application Configuration
    appConfiguration: frontend.Server {
       # Main Container Configuration
       mainContainer = container.Main {
10
            ports = [
                {containerPort = 80}
11
12
13
14-
       image = "nginx:1.7.8"
15
16
```

```
1 import base.pkg.kusion_models.kube.frontend
   import base.pkg.kusion_models.kube.frontend.service
   import base.pkg.kusion_models.kube.frontend.container
   import base.pkg.kusion models.kube.templates.resource as res tpl
   # Application Configuration
   appConfiguration: frontend.Server {
       # Main Container Configuration
       mainContainer = container.Main {
10
            ports = [
                {containerPort = 80}
11
12
13
14+
       image = "nginx:1.7.9"
15 }
16
```

Modules

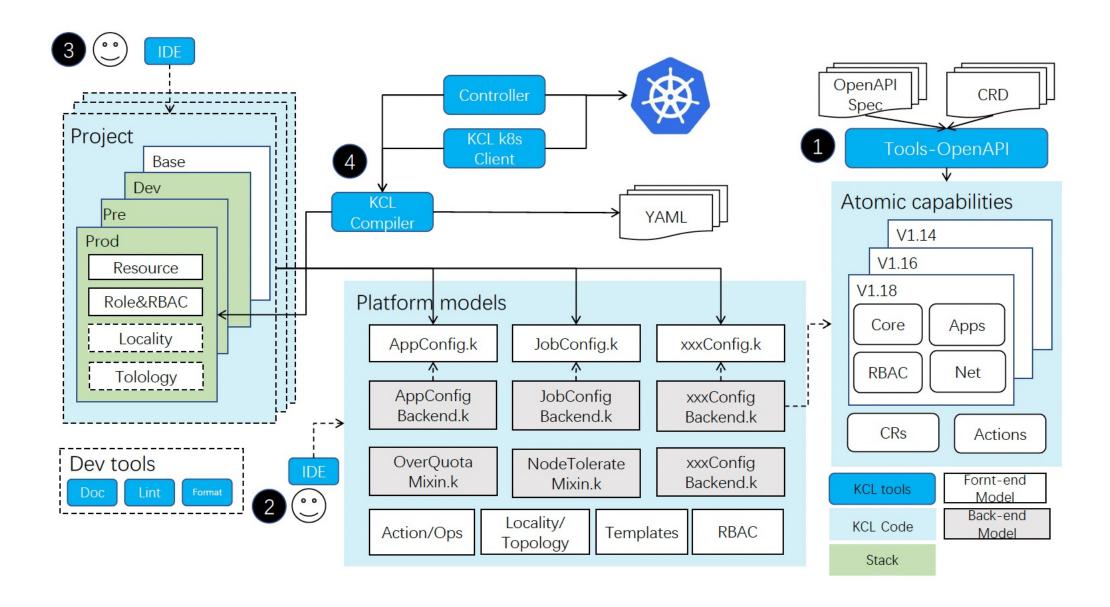




Workflow

03

Workflow





Evaluation

04

Related Works

Table 2: Features of configuration languages

			or coming ar					
	KCL	GCL	CUE	Jsonnet	HCL	Dhall		
Variables	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark		
Reference	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark		
Data types	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark		
Schema	\checkmark	✓	\checkmark	×	×	\checkmark		
Inherited	schema	tuple	×	obejct	data	data		
Arithmetic&Logic	\checkmark	✓	\checkmark	✓	\checkmark	\checkmark		
Loop	list/dict/schema	list	list	list/object	for splat	list generate function		
ьоор	comprehension	comprehension	comprehension	comprehension	(comprehension)	nst generate function		
Conditional	\checkmark	\checkmark	✓	\checkmark	\checkmark	\checkmark		
Built-in function	\checkmark	\checkmark	\checkmark	✓	\checkmark	\checkmark		
Function definition	✓	×	×	✓	×	\checkmark		
Import	\checkmark	✓	\checkmark	✓	\checkmark	\checkmark		
Type check	\checkmark	\checkmark	\checkmark	\times	×	\checkmark		
Testable	✓	×	×	✓	×	✓		
Mixin	\checkmark	×	×	×	×	×		
Data integration	✓	×	×	JSON	×	×		
Dynamic configuration	✓	×	×	✓	×	×		
Policy	\checkmark	×	×	\times	×	×		
v	idempotent merge/					idempetant marga/		
Merge	patch/	patch	idempotent merge	idempotent merg	idempotent merge	e patch	ch X	idempotent merge/ patch
	unique configuration					paten		

Key Results

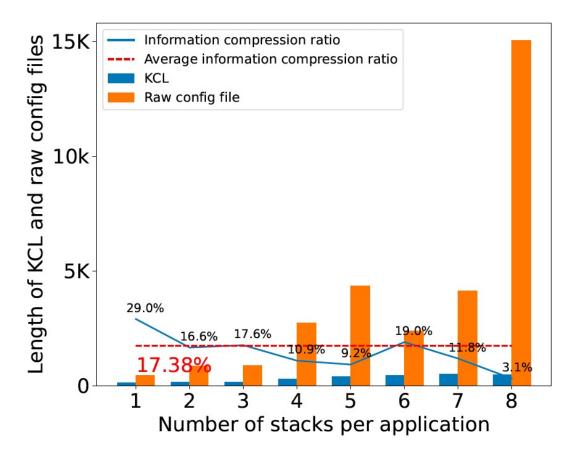


Fig. 10: Information compression ratio of KCL and raw config files

Key Results

10 K/day

~600 K

~18 K

KCL Compilations

KCL Code **PRs**

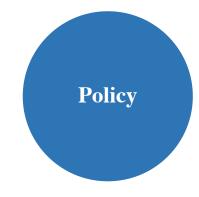
800

Projects

25 days – 2 days

Average configuration writing and shipping time

Future Work



• We will better support policy capabilities such as logic writing and data query to satisfy scenarios such as service authentication.



• We will provide a more complete KCL tool-set including create, query and update to meet more automation scenarios.



• We will improve language security through static model checking and theorem proving and let more problems be exposed to compile time as much as possible.

Summary

- We proposed the **KCL declarative language**, which is an open source constraint-based record & functional language mainly used in configuration and policy scenarios.
- In KCL, we are making special design including **modeling**, **constraint and workflow for the stability and scalability** of configuration, including configuration graph unification and carrying multiple configuration merging strategies.
- To date, the KCL has been used in more than **800 projects**, and the average configuration writing and distributing time is shortened from more than **25 days** to **2 days**. We have demonstrated the feasibility of our contribution through a large number of code practices and achievements.



THANKS